

REVERSIBLE EXECUTION AS A
DIAGNOSTIC TOOL
(Preliminary Draft)

by

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By the addition of only two new statements to the PL/I language, a programming system has been implemented which allows the user to retrace or backup execution of his program. He can backup N statements, to a label, to the last time that a variable was assigned a value, or until a given logical expression becomes true. Uses of this facility include the debugging of programs and applications where a stack is employed.

The first addition was a backup mechanism via a RETRACE statement which has the format RETRACE <CONDITION> AND <PL/I STATEMENT>, and has the effect that execution is retraced until <CONDITION> becomes true, the <PL/I STATEMENT> is executed, and then the program is re-executed from that point onwards.

The second change is a generalization of the PL/I ON statement by the inclusion of generalized logical expressions. The syntax of the ON statement was modified to be ON <EXPRESSION> <PL/I BEGIN BLOCK>. The user could activate an error routine by coding ON <EXPRESSION> CALL ERRORROUTINE, and whenever <EXPRESSION> was true, the ON unit would be activated.

With these two additions a powerful debugging facility emerges. After the error occurs the programmer can try to fix the source of the error rather than the effect. Less checks must be made while the program is executing in the forward direction. This is especially true in an interactive environment. At the very least the programmer could write a debugging routine that would "trace the N statements before an error is detected". Thus information local to the point of the error is produced without the production of the huge volumes of output typical with a tracing system.

Another use for this facility would be wherever a stack is used. Rather than including a push down stack within the program, the user could use the retrace mechanism to backup to the point before the entry was added in order to 'pop' the stack. This would be useful in the areas of symbol manipulation and compiler writing.

This system was implemented as part of the PL/I compiler called PL/C, and while it has not had customer usage, test results on randomly selected PL/C programs seem to show that performance is degraded by only a factor of two and that program size increased 40%. Execution times are significantly less than the usual times for interpretive debugging systems. Additional methods are displayed which can be used to microprogram this system in hardware. As hardware, the overhead should drop considerably, thus it would be possible to leave the debugging code within the program throughout its productive life in case an error should arise in the future.

The results show that such a facility is not very expensive to include as part of a compiler, and that future compilers would greatly benefit by the inclusion of such a debugging system. In many situations the power that such a debugging facility adds offsets the increase in execution time.

Example 1.

This program computes the roots of a quadratic equation; however, there is a keypunching error in statement 12 where an **B** appears instead of a **C**. The program executes until statement 12 which activates the error mechanism. This in turn activates the tracing mechanism which traces the remainder of the output.

Example 2.

The retrace statement has uses wherever a stack is employed. This example produces the left list of a tree by using the retrace statement to back up to a higher level in the tree when necessary. This type of routine is quite common in parsing algorithms such as in the first pass of a compiler. It is interesting since the backup mechanism is not being used solely for debugging in this case.

Example 3.

A binary search routine is exhibited which almost works since the end conditions are not tested properly. The backup mechanism insures that the program will continue processing with some reasonable assumptions being made. This enables the user to test all of his data, even if some of the values are incorrect. In this example, the second search (KEY =25) runs off the end of the table. The retrace insures that a reasonable value is used.

Example 4.

This example shows another non-debugging use of the retrace statement. Many times one wants to "undo" some code, such as removing an entry from a hash table. This is a difficult task; however, by use of the retrace statement to keep track of the various pointers, it becomes much easier to implement.

Please note that this system was implemented in version 4 of PL/C which did not include GR units; therefore, the logical conditional was implemented as an `IF` statement. Current plans are to install this system in version 5 where it will operate as described herein.

REFERENCES

1. Zelkowitz, H., Reversible Execution as a Diagnostic Tool, PhD Thesis, Cornell University, 1971
2. PL/C: The Cornell Compiler for PL/I, Department of Computer Science, Cornell University, 1976

AP 1.00000E+00 B= 4.00000E+00 C= 3.00000E+00
***** IN STMT 12 ERROR EXPT E HAS NOT BEEN INITIALIZED, IT IS SET TO 0.

#1(FU002) ACTIVE AT ST # 0012

ERRN UNCCNRD

#1(FU002) ACTIVE

AT ST # 0005 TO ST # 0011

B= 4.00000E+00

AT ST # 0012

***** IN STMT 12 ERROR EXPT E HAS NOT BEEN INITIALIZED, IT IS SET TO 0.

#1(FU002) ACTIVE AT ST # 0012

ERRN UNCCNRD

#1(FU002) ACTIVE

TRACE FAILS

B= 4.00000E+00

AT ST # 0012

***** IN STMT 12 ERROR EXPT E HAS NOT BEEN INITIALIZED, IT IS SET TO 0.

***** IN STMT 16 PROGRAM RETURNS FROM MAIN PROCEDURE

```

/* TEST PROGRAM 1 */
SYMI LEVEL NEST BLOCK SOURCE STATEMENT 19 FIELD

```

```

/* TEST PROGRAM 1 */
PI PROCEDURE OPTIONS(MAIN)

```

```

1
2
3
4
5
6
7
8
9
10
11
12
13
14

```

```

/* PROGRAM COMPUTES ROOTS OF EQUATION AX**2+BX+C=0 BY QUADRATIC */
/* FORMULA */
/* READ IN DATA */:R1 LIST(A,B,C)
/* PRINT DATA */:PUT DATA(A,B,C)
/* COMPUTE ROOTS */
K=0
DO I=1-2 SORTED=ASC(1/2)*A+
PUT SKIP LIST(ROOTS ARE:R001,R002)
/* THE DATA CARD CONTAINS THE VALUES 1,4,3 */
END
END P1

```

ERROR CDOC NO FILE SPECIFIED. SYMIN/SYSPRINT ASSUMED.

NEXT NODE IN TREE 1
NEXT NODE IN TREE 4
NEXT NODE IN TREE 3
NEXT NODE IN TREE 5

#RETRACE(0019) ACTIVE
AT ST # 0019 TO ST # 0011

BACKED UP TO NODE 3

#RETRACE(0019) ACTIVE
AT ST # 0019 TO ST # 0011

BACKED UP TO NODE 4
NEXT NODE IN TREE 2

#RETRACE(0019) ACTIVE
AT ST # 0019 TO ST # 0017

BACKED UP TO NODE 2

#RETRACE(0019) ACTIVE
AT ST # 0019 TO ST # 0011
IN STMT 23 PROGRAM IS STOPPED


```

1  /* PROGRAM 2 - PRINT A TREE IN LEFT LIST ORDER */
2  TREEPI PROCEDURE OPTIONS(MAIN);
3
4  DECLARE I TREE,
5          2 SON(20) FIXED BINARY,
6          2 BROTHER(20) FIXED BINARY;
7  /* READ IN TREE, EACH NODE CONTAINS A SON AND */
8  /* BROTHER POINTER */
9  GET LISTNUMBER OF NODES11
10
11  /* READ IN TREE */
12  /* ROOT OF TREE IS AT NODE 1 */
13  J=1; PUSH1;
14  /* AT NODE J, MOVE TO SON OF J */
15  L1=SON(J);
16  /* IF NOT AT LEAF CONTINUE DOWN STRUCTURE */
17  IF 120 THEN DO;
18    PUSH2; J=1; GO TO L11 END;
19  /* GO TO BROTHER */
20  L2; J=BROTHER(J);
21  /* IF NOT LEAF CONTINUE */
22  IF J20 THEN DO;
23    PUSH3; GO TO L1; END;
24  /* BACKUP TO FATHER POINTER */
25  /* SEE IF FINISHED */
26  IF J=1 THEN STOP;
27  /* SET TO READ THIS OPTION NEXT TIME RETAGE IS CALLED */
28  DEBUG=0;
29  /* PRINT CURRENT NODE */
30  PUT SKIP LIST( ' ', ' ', BACKED UP TO NODE', J11
31  GO TO L2;
32
33  /* DATA IS 5 4 0 0 0 5 0 3 2 0 0 */
34  /* WHICH IS TREE (1) (2) (3) (4) (5) (6) (7) (8) (9) (10) */
35  -<2>
36  END TREEPI;
  
```

***** IN STMT 15 ERRORS EXRT N HAS NOT BEEN INITIALIZED. IT IS SET TO 0.
***** IN STMT 15 ERRORS EXRT NEXT_ITEM HAS NOT BEEN INITIALIZED. IT IS SET TO 0.

KEY= 4

WIP(0000) ACTIVE AT ST # 0039

WCLPACT(0000) ACTIVE

AT ST # 0005 TO ST # 0037

KEY= 15

KEY= -1

1 0
2 1
3 1
4 0
5 0
6 0
7 0
8 0
9 0
10 0
11 0
12 0
13 0
14 0
15 1
16 0
17 0
18 0
19 1
20 0
21 0
IN STMT 22 PROGRAM IS STOPPED

```

/* TEST PROGRAM 3 */
SUBJECT PROCEDURE OPTIONS(TABLE)
/* THIS PROCEDURE DOES A BINARY SEARCH OF ITEMS IN A TABLE */
/* WHEN FOUND IT INCREMENTS A COUNTER. THIS SHOULD BE USEFUL */
/* IN A SYSTEM THAT KEPT TRACK OF THE NUMBER OF ACCESSES */

2 /* DECLARE VARIABLES */
   DECLARE I TABLE(20),
           I KEY, 2 ACCESS 1 FIXED BINARY,
           DIFFERENCE FIXED BINARY;

3 /* SET LEAVE EXIT IF SUBSCRIPT RUNS OFF TABLE */
   IF NEXT_ITEM# THEN GO;
4   /* TRACE (DEBUG) CONDITION(NEXT_ITEM<N) AND CALL FNDI ENDI
5   /* NEXT_ITEM# SO THEN DOI
6   /* (NEXT_ITEM#<0) CONDITION(NEXT_ITEM#>0) AND CALL FNDI ENDI

15 ACCESS=0;
16 /* READ IN TABLE */
   /* DATA READ IN IS: 9 1 9 7 10 10 10 19 21 */
17 /* READ IN DESIRED LOCATION */
   LOOP: GET LISTEN(KEY); PUT SKIP EDIT(NEW_KEY,NEW_KEY)
18   /* QUIT IF FINISHED */ IF NEW_KEY<0 THEN DO;
19     PUT EDIT(KEY,ACCESS) DO 1=1 TO N11;
20     /* SEARCH TABLE */
21     /* (SKIP 2 F33); STOP 5 END;
22     /* SEARCH TABLE */
23     /* SEARCH TABLE */
24     /* SEARCH TABLE */
25     /* SEARCH TABLE */
26     /* SEARCH TABLE */
27     /* SEARCH TABLE */
28     /* SEARCH TABLE */
29     /* SEARCH TABLE */
30     /* SEARCH TABLE */

31 /* FOUND ITEM IN LIST */
32   /* FOUND: ACCESS(NEXT_ITEM)=ACCESS(NEXT_ITEM)+1;
33     PUT LIST(ACCESS(NEXT_ITEM)); GO TO LOOP;
34   /* TOO FAR IN TABLE, NO COPY */
35   /* UPPEK: DIFFERENCE=(DIFFERENCE+1)/2;
36     NEXT_ITEM=NEXT_ITEM-DIFFERENCE; GO TO SRCH1
37   /* NOT FAR ENOUGH IN TABLE */
38   /* DOWNK: DIFFERENCE=(DIFFERENCE+1)/2;
39     NEXT_ITEM=NEXT_ITEM+DIFFERENCE; GO TO SRCH1

40 /* READ TRACE OPTION NEXT TIME */
41 FNDI PROCEDURE;
42 DEPEND=0; GO TO FNDI; END FNDI;
43
44 END UNSPCH;

```

ERROR -
#REF(10018) ACTIVE 4
AT ST # 0018 TO ST # 0006
ERROR 6
#REF(10018) ACTIVE
AT ST # 0018 TO ST # 0006
TABLE CONTENTS
1 2 3 4 5 6 7 8 9 10
IN STMT 11 PROGRAM IS STOPPED

PL/C (DEBUG)

-C20./- TEST PROGRAM 4 */

PL/C-4420013 21/07/51 11172 PAGE

SYMT LEVEL NEXT BLOCK

SOURCE STATEMENT

10 FIELD

```
1 /* TEST PROGRAM 4 */
2 /* PROCEDURE OPTIO(SIMAINI)
3 /* THIS IS AN EXAMPLE WHERE RETRACE IS USED TO UNDO A COMPLICATED */
4 /* SECTION OF CODF, SUCH AS THE DELETION OF ENTRIES IN A HASH TABLE*/
5 DECLARE TABLE(S) FIXED BINARY, TOTAL FIXED BINARY, A(I) FIXED BINARY
6 TOTAL=0; IND=0;
7 /* READ INPUT DATA */
8 GET LIST(I) NO 1 TO 11111
9 /* READ NEXT ENTRY INTO N */
10 LOOP: IND=IND+1; N=A(IND);
11 /* IF N IS NEGATIVE, QUIT */
12 IF N<0 THEN DO;
13 PUT SKIP EDIT('TABLE CONTENTS: (TABLE) NO. 1
14 1 TO TOTAL(A),SKIP, (TOTAL) F5)11;
15 STOP; END;
16 /* INSERT N INTO TABLE (SUCH AS HASH TABLE) */
17 CALL INSERT(N);
18 /* COMPUTE SOME FUNCTION BASED UPON N */
19 B=F(N);
20 /* DELETE ENTRY IF B FAILS SOME CONDITION */
21 IF (B=1) (B=36) THEN DO;
22 PUT SKIP LIST('ERROR--');
23 RETRACE(DEBUG) TO (LOOP) AND CALL BACKUP; END;
24 GO TO LOOP;
25 /* BACKUP PROCEDURE CALL */
26 BACKUP: PROCEDURE;
27 /* SET DEBUG OPTION */
28 DEBUG=0;
29 /* REREAD OLD DATA */
30 IND=IND+1;
31 /* GET NEXT DATA ITEM */
32 GO TO LOOP; END BACKUP;
33 /*INSERT ROUTINE - THIS CAN BE AS COMPLICATED AS DESIRED */
34 INSERT:PROCEDURE; TOTAL=TOTAL+1; TABLE(TOTAL)=N; END INSERT;
35 /* FUNCTION COMPUTATION */
36 F1:PROCEDURE; RETURN(N); END F1;
37 /* DATA CARDS ARE 1 2 3 4 5 6 7 8 9 10 -1 */
38 END P;
```